CSE 413, Spring 2011, Assignment 4
Due: Tuesday 3 May, 11:00PM

Set-up: For this assignment, edit a copy of `hw4skeleton.scm`, which is on the course website. In particular, replace occurrences of `"CHANGE"` to complete the problems. Do not use any mutation (`set!`, `set-car!`, etc.) anywhere in the assignment.

Overview: This homework has to do with MUPL (a Made Up Programming Language). MUPL programs are written directly in Scheme (using Pretty Big in DrRacket) using the structs defined at the beginning of `hw4skeleton.scm`, according to this syntax definition:

- If s is a Scheme string, then `(make-var` s`)` is a MUPL expression (a variable use).
- If n is a Scheme integer, then `(make-int` n`)` is a MUPL expression (a constant).
- If e_1 and e_2 are MUPL expressions, then `(make-add` e_1 e_2`)` is a MUPL expression (an addition).
- If s_1 and s_2 are Scheme strings and e is a MUPL expression, then `(make-fun` s_1 s_2 e`)` is a MUPL expression (a function). In e, s_1 is bound to the function itself (for recursion) and s_2 is bound to the (one) argument. Also, `(make-fun #f` s_2 e`)` is allowed for nonrecursive functions.
- If e_1, e_2, and e_3, and e_4 are MUPL expressions, then `(make-ifgreater` e_1 e_2 e_3 e_4`)` is a MUPL expression (a conditional meaning e_1 is strictly greater than e_2).
- If e_1 and e_2 are MUPL expressions, then `(make-app` e_1 e_2`)` is a MUPL expression (a function application).
- If s is a Scheme string and e_1 and e_2 are MUPL expressions, then `(make-mlet` s e_1 e_2`)` is a MUPL expression (a let expression) where the value of e_1 is bound to s in the evaluation of e_2.
- If e_1 and e_2 are MUPL expressions, then `(make-apair` e_1 e_2`)` is a MUPL expression (a pair-creator).
- If e_1 is a MUPL expression, then `(make-fst` e_1`)` is a MUPL expression (getting part of a pair).
- If e_1 is a MUPL expression, then `(make-snd` e_1`)` is a MUPL expression (getting part of a pair).
- `(make-aunit)` is a MUPL expression (holding no data, much like `()` in ML or Scheme).
- If e_1 is a MUPL expression, then `(isaunit` e_1`)` is a MUPL expression (testing for `(make-aunit)`).
- `(make-closure` *env* f`)` is a MUPL value where f is MUPL function (an expression made from `make-fun`) and *env* is an environment mapping variables to values. Closures do not appear in source programs; they result from evaluating functions.

A MUPL *value* is an integer constant (wrapped in `make-int`), a closure, unit, or a pair of values. Notice that like in Scheme we can build list values out of nested pair values that end with unit.
You should assume MUPL programs are syntactically correct (e.g., do not worry about wrong things like `(make-int "hi")` or `(make-int (make-int 37))`. But do *not* assume MUPL programs are free of "type" errors like `(make-add (make-aunit) (make-int 7))` or `(make-fst (make-int 7))`.

Warning: This assignment is difficult because you have to understand MUPL well and debugging an interpreter is an acquired skill. Start early.

1. (Implementing the language) Write a MUPL interpreter, i.e., a Scheme function `eval-prog` that takes a MUPL program `p` and either returns the MUPL value that `p` evaluates to or calls Scheme's `error` if evaluation encounters a run-time MUPL type error or unbound MUPL variable.

   A MUPL expression is evaluated under an environment (for evaluating variables, as usual). Use a list of pairs for the environment (starting with `()`) so that you can use the provided `envlookup` function. Here is an informal semantics for MUPL expressions:

- All values (including closures) evaluate to themselves. For example, `(eval-prog (make-int 17))` would return `(make-int 17)`, *not* 17.
- A variable evaluates to a value according to the environment.
- An addition evaluates its subexpressions and assuming they both produce integers, produces the integer that is their sum. (Note this case is done for you to get you pointed in the right direction.)
- Functions are lexically scoped: A function evaluates to a closure holding the function and the current environment.
- An `ifgreater` evaluates its first two subexpressions to values v_1 and v_2 respectively. Assuming both values are integers, it evaluates its third subexpression if v_1 is a strictly greater integer than v_2 else it evaluates its fourth subexpression.
- An application evaluates its first and second subexpressions to values. If the first is not a closure, it is an error. Else, it evaluates the closure's function's body in the closure's environment extended to map the function's name to the closure (unless the name field is `#f`) and the function's argument to the second value of the application.
- An mlet expression evaluates its first expression to a value v. Then it evaluates the second expression to a value, in an environment extended to map the name in the mlet expression to v.
- A pair expression evaluates its two subexpressions and produces a (new) pair holding the results.
- A fst expression evaluates its subpexpression. It is an error if the result is not a pair of values. Else the result of the fst expression is the `e1` field in the pair.
- A snd expression is the same as a fst expression except the result is the `e2` field of the pair.
- An `isaunit` expression evaluates its subpexpression. If the result is unit, then the result for the isunit expression is the integer 1, else the result is the integer 0.

Hint: The `app` case is definitely the most complicated. In the sample solution, no case is more than 12 lines and several are 1 line.

2. (Expanding the language) MUPL is a small language, but we can write Scheme functions that act like MUPL macros. They produce MUPL programs that other code could later pass to `eval-prog`. These Scheme functions you write produce MUPL expressions that you could have written by hand, i.e., they are "macros for MUPL." In implementing these Scheme functions, do not use `make-closure` (which is only used internally in `eval-prog`) nor `eval-prog` (we are creating a program, not running it).

   (a) Write a Scheme function `ifunit` that takes three MUPL expressions e_1, e_2, and e_3. It returns a MUPL expression that when run evaluates e_1 and if the result is unit then it evaluates e_2 and that is the overall result, else it evaluates e_3 and that is the overall result. Sample solution: 1 line.

   (b) Write a Scheme function `mlet*` that takes a list of lists $((s_1, e_1) \ldots (s_i, e_i) \ldots (s_n, e_n))$ and a final MUPL expression e_{n+1}. In each pair (s_i, e_i), the s_i is a string and e_i is a MUPL expression. `mlet*` returns a MUPL expression whose value is e_{n+1} evaluated in an environment where each s_i is bound to the result of evaluating the corresponding e_i for $1 \leq i \leq n$. The bindings are done sequentially, so that each e_i is evaluated in an environment where s_1 through s_{i-1} have been previously bound to the values e_1 through e_{i-1}.

   (c) Write a Scheme function `ifeq` that takes four MUPL expressions e_1, e_2, e_3, and e_4 and returns a MUPL expression that acts like `ifgreater` except e_3 is evaluated if e_1 and e_2 are equal integers. Unfortunately, MUPL does not have hygiene and we want to evaluate e_1 and e_2 exactly once, so assume the MUPL expressions do not use the variables `_x` and `_y` (i.e., you can use these variables to implement `ifeq`). Sample solution is 7 (short) lines.

3. (Using the language) We can write MUPL expressions directly in Scheme using the constructors for the structs and (for convenience) the functions we wrote in the previous problem.

(a) Bind to the Scheme variable `mupl-map` a MUPL function that acts like map (as we used in Scheme). Your function should be curried: it should take a MUPL function and return a MUPL function that takes a MUPL list and applies the function to every element of the list returning a new MUPL list. A MUPL list is simply unit or a pair where the second component is a MUPL list. Sample solution: 7 lines.

(b) Bind to the Scheme variable `mupl-mapAddN` a MUPL function that takes an integer i and returns a MUPL function that takes a list of integers and returns a new list that adds i to every element of the list. Use `mupl-map` (a use of `mlet` is given to you to make this easy). Sample solution is 4 lines (including the line given to you).

4. **Challenge Problem:** Write a second version of `eval-prog` (bound to `eval-prog2`) that builds closures with smaller environments: When building a closure, it uses an environment that is like the current environment but only holds variables that are free variables in the function part of the closure. Note: You will have to write a Scheme function that takes a MUPL expression and computes its free variables.

   For full challenge-problem credit, use memoization (yes, you should use mutation for this) to avoid computing any function's free variables more than once.

   Warning: The sample solution does *not* include a solution to the extra credit.

Turn-in Instructions

- Put all your solutions in one file, `hw4.scm` and turn in that file using the regular online dropbox.
- The first line of your `.scm` file should be a Scheme comment with your name and the phrase `CSE 413, Spring 2011, Homework 4`.